

Zeno Effect for Quantum Computation and Control

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(Received 28 April 2011; revised manuscript received 9 February 2012; published 22 February 2012)

It is well known that the quantum Zeno effect can protect specific quantum states from decoherence by using projective measurements. Here we combine the theory of weak measurements with stabilizer quantum error correction and detection codes. We derive rigorous performance bounds which demonstrate that the Zeno effect can be used to protect appropriately encoded arbitrary states to arbitrary accuracy while at the same time allowing for universal quantum computation or quantum control.

DOI: 10.1103/PhysRevLett.108.080501

PACS numbers: 03.67.Pp, 03.65.Xp, 03.65.Yz

Protection of quantum states or subspaces of open systems from decoherence is essential for robust quantum information processing and quantum control. The fact that measurements can slow down decoherence is well known as the quantum Zeno effect (QZE) [1,2] (for a recent review, see Ref. [3]). The standard approach to the QZE uses repeated strong, projective measurements of some observable V . In this setting, it can be shown that such repeated measurements decouple the system from the environment or bath and project it into an eigenstate or eigensubspace of V [4,5]. Projective measurements are, however, an idealization, and here we are interested in the more realistic setting of weak, nonselective measurements, implementing a weak-measurement quantum Zeno effect (WMQZE). The measurements are called weak since all outcomes result in small changes to the state [6,7] and nonselective since the outcomes are not recorded. Such measurements can include many phenomena not captured by projective measurements, e.g., detectors with nonunit efficiency, measurement outcomes that include additional randomness, and measurements that give incomplete information [see Eq. (1)]. The WMQZE has already been considered in a wide range of applications, e.g., Refs. [8–11]. However, a general systematic study of decoherence suppression via the WMQZE, allowing for universal quantum control, appears to be lacking. This work aims at bridging this gap. More specifically, we ask whether the WMQZE can be used to protect arbitrary quantum states while they are being controlled, e.g., for the purpose of quantum computation. Borrowing quantum coding ideas, we devise a measurement protocol which allows us to provide an affirmative answer to this question: Namely, assume that the system-bath interaction is local and bounded and that we can encode arbitrary system states into a sufficiently large stabilizer quantum error correcting code. Then a weak system-measurement protocol using M stabilizer measurements of strength ϵ lasting a

total time τ suppresses the system-bath interaction arbitrarily well in the limit of large M and commutes with quantum control or computation being performed on the system. We dedicate the rest of this work to explaining, sharpening, and proving this claim.

Weak measurements.—A generalized, positive operator-valued measure comprises a set of “measurement operators” $\{M_j\}$ satisfying the sum rule $\sum_j M_j^\dagger M_j = \mathbb{1}$, which map a state ρ to $\rho_j = M_j \rho M_j^\dagger / p_j$ with probability $p_j = \text{Tr}[M_j \rho M_j^\dagger]$ for measurement outcome j [12]. In general, one can write the weak-measurement superoperator corresponding to a two-outcome measurement of an observable (Hermitian operator) V with strength ϵ on a state ρ as [13] $\mathcal{P}_\epsilon(\rho) = \sum_{r=\pm} P_V(r\epsilon) \rho P_V(r\epsilon)$, where $P_V(\epsilon) = \sum_{s=\pm} \alpha_s(\epsilon) P_{sV}$, with $P_{\pm V} \equiv \frac{1}{2}(\mathbb{1} \pm V)$ standard projection operators when $V^2 = \mathbb{1}$, and $\alpha_\pm(\epsilon) \equiv \sqrt{[1 \pm \tanh(\epsilon)]/2}$. Since $P_V^2(\epsilon) + P_V^2(-\epsilon) = \mathbb{1}$, the operators $\{P_V(\epsilon), P_V(-\epsilon)\}$ satisfy the sum rule and, hence, are measurement operators for a given ϵ . They are parametrized by the strength ϵ so that they can be considered weak-measurement operators. Since $\lim_{\epsilon \rightarrow \pm\infty} P_V(\epsilon) = P_{\pm V}$, the ideal or strong measurement limit is recovered when the measurement strength $|\epsilon| \rightarrow \infty$, i.e., $\mathcal{P}_\infty(\rho) = \sum_{s=\pm} P_{sV} \rho P_{sV}$. The no-measurement scenario is the case $\epsilon \rightarrow 0$, i.e., $\mathcal{P}_0(\rho) = \rho$. The weak measurement of an operator V with strength ϵ can be rewritten as

$$\mathcal{P}_\epsilon(\rho) = (1 - \zeta) \mathcal{P}_\infty(\rho) + \zeta \rho, \quad \zeta \equiv \text{sech}(\epsilon), \quad (1)$$

and thus a weak measurement can be interpreted as a noisy measurement in which, with probability ζ , the measurement is not executed [14]. A strong measurement is the idealized case, when $\zeta = 0$. Weak measurements are universal in the sense that they can be used to build up arbitrary measurements without the use of ancillas [13].

Open system evolution with measurements.—Consider a system and bath with respective Hilbert spaces \mathcal{H}_S and

\mathcal{H}_B . The joint evolution is governed by the Hamiltonian $H = H_0 + H_{SB}$, where $H_0 \equiv H_S \otimes \mathbb{1}_B + \mathbb{1}_S \otimes H_B$, acting on the joint Hilbert space $\mathcal{H}_{SB} \equiv \mathcal{H}_S \otimes \mathcal{H}_B$. We assume that $\|H_\mu\| \equiv J_\mu/2 < \infty$ ($\mu \in \{0, S, B, SB\}$) [15]. We denote $J_1 \equiv J_{SB}$. Thus $\|H\| \leq (J_0 + J_1)/2 \equiv J/2 < \infty$.

We wish to protect an arbitrary and unknown system state ϱ_S against decoherence for some time τ by using only weak measurements. We model all such measurements as instantaneous and perform M equally spaced measurements in the total time τ . We define superoperator generators $\mathcal{L}_\mu(\cdot) \equiv -i[H_\mu, \cdot]$ and $\mathcal{L}(\cdot) \equiv -i[H, \cdot]$. The free evolution superoperator $\mathcal{U}(\tau) \equiv e^{\mathcal{L}\tau}$ describes the evolution after each measurement. Hence, the joint state and system-only state, after time τ , are given by

$$\varrho_{SB}(\tau) = \left[\mathcal{P}_\epsilon \mathcal{U}\left(\frac{\tau}{M}\right) \right]^M \varrho_{SB}(0), \quad \varrho_S(\tau) = \text{Tr}_B \varrho_{SB}(\tau). \quad (2)$$

From now on we shall assume for simplicity that the initial system state is pure: $\varrho_S(0) = |\psi_S(0)\rangle\langle\psi_S(0)|$ and that the joint initial state is factorized, i.e., $\varrho_{SB} = \varrho_S \otimes \varrho_B$. For notational simplicity we denoted $\varrho_\mu(0) \equiv \varrho_\mu$. Note that in Eq. (2) \mathcal{P}_ϵ acts nontrivially only on system operators.

Figure of merit.—To determine the success of our protection protocol, we compare the “real” system state with protection and in the presence of H_{SB} [Eq. (2)] to the uncoupled ($H_{SB} = 0$), unprotected “ideal” system state, namely, to $\varrho_S^0(\tau) = \text{Tr}_B \varrho_{SB}^0(\tau)$, with $\varrho_{SB}^0(\tau) = \mathcal{U}_0(\tau) \varrho_{SB}$, where $\mathcal{U}_0(\tau) \equiv e^{\mathcal{L}_0 \tau}$ and $\mathcal{L}_0 = \mathcal{L}_S + \mathcal{L}_B$ ($[\mathcal{L}_S, \mathcal{L}_B] = 0$) are the ideal unitary superoperator and its generator, respectively.

A suitable figure of merit is then the trace-norm distance [12,15] $D[\varrho_1, \varrho_2] \equiv \frac{1}{2} \|\varrho_1 - \varrho_2\|_1$ between the real and ideal states. We shall show that we can make $D[\varrho_S(\tau), \varrho_S^0(\tau)]$ arbitrarily small for a given H by a suitable choice of weak measurements.

Weak measurements over a stabilizer code.—Previous WMQZE work applied only to particular states [8–11]. To achieve our goal of protecting an arbitrary, unknown k -qubit state, we encode the state into an $[[n, k, d]]$ stabilizer quantum error correcting code (QECC) [12,16], with stabilizer group $\mathbf{S} = \{S_i\}_{i=0}^{\bar{Q}}$, where $S_0 \equiv \mathbb{1}$. We assume that the code distance $d \geq 2$; i.e., the code is at least error-detecting, with generators $\bar{\mathbf{S}} = \{\bar{S}_i\}_{i=1}^{\bar{Q}} \subset \mathbf{S}$, where $\bar{Q} = n - k$. Note that every stabilizer element can be written as $S_i = \prod_{\nu=1}^{\bar{Q}} \bar{S}_\nu^{r_{i\nu}}$, where $r_{i\nu} \in \{0, 1\}$; i.e., the stabilizer elements are given by all possible products of the generators, whence $Q + 1 = 2^{\bar{Q}}$. The encoded initial state $|\psi_S(0)\rangle$ is a simultaneous $+1$ eigenstate of all the elements of \mathbf{S} . We can associate a pair of projectors (measurement operators) $P_{\pm S_i} \equiv \frac{1}{2}(\mathbb{1} \pm S_i)$ to each stabilizer-group element and, accordingly, a pair of weak-measurement operators $\{P_{S_i}(\epsilon), P_{S_i}(-\epsilon)\}$ to each S_i , i.e., $P_{S_i}(\epsilon) = \sum_{s=\pm} \alpha_s(\epsilon) P_{s S_i}$. In quantum error correction (QEC) one performs a strong measurement of the generators in order

to extract an error syndrome [16]. It has been recognized that these strong syndrome measurements implement a QZE [17,18]. When we measure $\bar{\mathbf{S}}$, we need to form products of the weak-measurement operators of all the generators, accounting for all possible sign combinations. Let $P_{\bar{\mathbf{S}}}^{(b)}(\epsilon) \equiv \prod_{i=1}^{\bar{Q}} P_{\bar{S}_i}^{[(-1)^{b_i} \epsilon]}$ denote such a product for a given choice of signs uniquely determined by the integer $b = \sum_{i=0}^{\bar{Q}} b_i 2^i$, with $b_i \in \{0, 1\}$. Letting

$$\bar{\mathcal{P}}_\epsilon(\varrho) = \sum_{b=0}^{2^{\bar{Q}}-1} P_{\bar{\mathbf{S}}}^{(b)}(\epsilon) \varrho P_{\bar{\mathbf{S}}}^{(b)}(\epsilon), \quad (3)$$

we can now define a weak stabilizer generator measurement protocol as $[\bar{\mathcal{P}}_\epsilon \mathcal{U}(\frac{\tau}{M})]^M$.

We stress the two important differences between this protocol and the analogous stabilizer measurement step in QEC: First, we do not need to observe or use the syndrome; second, we allow for weak measurements. In this sense our assumptions are weaker than those of QEC, and hence the ability to perform QEC implies the ability to perform our protocol.

Moreover, for the same reason that the many-body character of stabilizer measurements is not a significant drawback in QEC theory, it is not a problem for our protocol either. The reason is that such measurements can be implemented (even fault-tolerantly) by using at most two-local operations. See Ref. [14] for an explicit two-local construction for the weak-measurement case. An alternative is to consider a protocol based on measuring the gauge operators of the Bacon-Shor code [19], which are all two-local and can be shown to implement a WMQZE as well [20].

We shall also consider a weak stabilizer-group measurement protocol: $[\mathcal{P}_\epsilon \mathcal{U}(\frac{\tau}{M})]^M$, where $\mathcal{P}_\epsilon(\varrho) = \sum_{b=0}^{2^{\bar{Q}}-1} P_{\bar{\mathbf{S}}}^{(b)}(\epsilon) \varrho P_{\bar{\mathbf{S}}}^{(b)}(\epsilon)$, with $P_{\bar{\mathbf{S}}}^{(b)}(\epsilon) \equiv \prod_{i=0}^{\bar{Q}} P_{S_i}^{[(-1)^{b_i} \epsilon]}$ and $b = \sum_{i=0}^{\bar{Q}} b_i 2^i$. As we shall see, the generators and group protocols exhibit substantial trade-offs, so we shall consider both in our general development below.

Note that if ϱ_S is stabilized by $\bar{\mathbf{S}}$ (or \mathbf{S}), then the weak-measurement protocol perfectly preserves an arbitrary encoded state in the absence of system-bath coupling. Another important fact we shall need later is that, given some $[[n, k, d]]$ stabilizer QECC, if a Pauli group operator P anticommutes with at least one of the stabilizer generators, then it anticommutes with half of all the elements of the corresponding stabilizer group \mathbf{S} [14].

Distance bound.—Following standard conventions, we call a Pauli operator k -local if it contains a tensor product of k nonidentity Pauli operators. We call a system Hamiltonian k -local if it is a sum of k -local Pauli operators, and a system-bath Hamiltonian k -local if it is a sum of k -local Pauli operators acting on the system, tensored with arbitrary bath operators.

Let $\vec{r}_s = \{r_i\}_{i=1}^{\bar{Q}}$, where $r_i \in \{0, 1\} \forall i$ and $s \in \{1, \dots, \bar{Q}\}$. Let $\{\Omega_i\}_{i=1}^{\bar{Q}}$, where $\Omega_i^2 = \mathbb{1} \forall i$, denote a commuting set

of operators acting on the system only. Consider the recursive definition $H_{\tilde{r}_s} = \frac{1}{2}[H_{\tilde{r}_{s-1}} + (-1)^s \Omega_s H_{\tilde{r}_{s-1}} \Omega_s]$, where $H_{\tilde{r}_0} \equiv H$. This construction allows for the decomposition of any Hamiltonian as $H = \sum_{\tilde{r}_Q} H_{\tilde{r}_Q}$, with the property $\{H_{\tilde{r}_Q}, \Omega_i\} = 0$ if $r_i = 1$ or $[H_{\tilde{r}_Q}, \Omega_i] = 0$ if $r_i = 0$. Note that $H_0 = H_{\tilde{r}_0}$ and $H_{SB} = \sum_{\tilde{r}_Q} H_{\tilde{r}_Q} - H_0$. It follows from the triangle inequality, norm submultiplicativity, and the recursive definition of $H_{\tilde{r}_i}$ that $\|H_{\tilde{r}_Q}\| \leq J_1/2$. These bounds can be further specialized or tightened for specific forms of the Hamiltonian.

We are now ready to state our main result.

Theorem 1.—Assume an arbitrary pure state $\varrho_S = |\psi_S\rangle\langle\psi_S|$ is encoded into an $[[n, k, d]]$ stabilizer QECC. Assume that $H_{SB} = \sum_{K=1}^{d-1} H_{SB}^{(K)}$ and that H_S commutes with the code's stabilizer, so that $H_S = \sum_{l \geq 1} H_S^{(ld)}$, where $H_{SB}^{(K)}$ ($H_S^{(K)}$) denotes a K -local system-bath (system-only) Hamiltonian, and all Hamiltonians, including H_B , are

bounded in the sup-operator norm. Finally, let $Q = 2^{n-k} - 1$ and $q = (Q + 1)/2$, and assume $J_0 > J_1$ (the case $J_0 < J_1$ will be discussed in Ref. [20]). Then the stabilizer-group measurement protocol $[\mathcal{P}_\epsilon \mathcal{U}(\tau/M)]^M$ protects ϱ_S up to a deviation that converges to 0 in the large M limit:

$$D[\varrho_S(\tau), \varrho_S^0(\tau)] \leq \left(1 - \frac{\Gamma_g(M)}{\Gamma_\parallel(M)}\right) [1 + \Gamma_\parallel(M)]^M - e^{\tau J_0} + \Gamma_g(M) \sum_{s=\pm} A_s(M) \gamma_s^{M-1}(M) \equiv B \quad (4a)$$

$$= e^{\tau J_0} \left[\frac{Q\tau^2 J_1^2}{4} + \frac{Q\tau J_1}{2} (1 + \tau J_0) \frac{\zeta^q}{1 - \zeta^q} \right] \times \frac{1}{M} + O\left(\frac{1}{M^2}\right), \quad (4b)$$

where

$$\Gamma_\parallel(M) := \frac{1}{Q+1} e^{\tau J_0/M} (e^{\tau Q J_1/M} + Q e^{-\tau J_1/M}) - 1, \quad \Gamma_g(M) := \frac{1}{Q+1} e^{\tau J_0/M} (e^{\tau Q J_1/M} - e^{-\tau J_1/M}), \quad (5a)$$

$$\gamma_\pm(M) := \frac{1}{2} [1 + \Gamma_\parallel + (1 + Q\Gamma_\parallel)\zeta^q] \pm \frac{1}{2} \sqrt{[1 + \Gamma_\parallel - (1 + Q\Gamma_\parallel)\zeta^q]^2 + 4Q\Gamma_\parallel^2 \zeta^q}, \quad (5b)$$

$$A_\pm(M) := \frac{Q\Gamma_\parallel \zeta^q (\gamma_\pm + \Gamma_\parallel) + (1 + \Gamma_\parallel) [(1 + \Gamma_\parallel) - \gamma_\mp]}{\Gamma_\parallel (\gamma_\pm - \gamma_\mp)}. \quad (5c)$$

For a generator measurement protocol $[\tilde{\mathcal{P}}_\epsilon \mathcal{U}(\tau/M)]^M$, replace q by 1 in Eqs. (4b), (5b), and (5c). In the strong measurement limit ($\epsilon \rightarrow \infty$), both protocols yield the distance bound

$$D[\varrho_S(\tau), \varrho_S^0(\tau)] \leq e^{J_0 \tau} \left[\left(\frac{Q e^{-J_1 \tau/M} + e^{J_1 \tau Q/M}}{Q+1} \right)^M - 1 \right]. \quad (6)$$

To motivate the locality aspects of Theorem 1, recall that by construction of a stabilizer code any Pauli operator with locality $\leq d - 1$ anticommutes with at least one stabilizer generator, a condition satisfied by all $H_{SB}^{(K)}$ in Theorem 1. Moreover, logical operators of the code (elements of the normalizer, which commute with the stabilizer) must have locality that is an integer multiple of the code distance d , a condition satisfied by every $H_S^{(ld)}$, which by assumption can be used to implement logical operations on the code while stabilizer measurements are taking place. To keep the locality of H_S low thus requires a low distance code. We present an example of a $d = 2$ code below.

Proof sketch of Theorem 1.—We first consider the case of weak measurements of the entire stabilizer group, \mathcal{P}_ϵ . A typical K -local term we need to calculate is then of the form $\mathcal{P}_\epsilon(H_{SB}^{(K)} \varrho_{SB}) = \sum_{b=0}^{2^Q-1} P_S^{(b)}(\epsilon) H_{SB}^{(K)} \varrho_{SB} P_S^{(b)}(\epsilon)$. Now we use the previously established fact that if E (modulo logical operators and stabilizer operations) is a correctable

error, then $\{S_i, E\} = 0$ for exactly half of the stabilizer elements. Hence the same number of stabilizer elements, $q = (Q + 1)/2$, anticommute with $H_{SB}^{(K)}$. From here a straightforward calculation reveals that $(\mathcal{P}_\epsilon)^j \mathcal{L}_{SB}(\varrho_{SB}) = \zeta^{jq} \mathcal{L}_{SB}(\varrho_{SB})$, a key result since it shows how the measurements suppress the “erred” component of the state, $\mathcal{L}_{SB}(\varrho_{SB})$. On the other hand, since we assume that $[H_S, S_i] = 0$ for all stabilizer elements, we have $\mathcal{P}_\epsilon[\mathcal{L}_0(\varrho_{SB})] = \mathcal{L}_0(\varrho_{SB})$ and hence $(\mathcal{P}_\epsilon)^j[\mathcal{L}_0(\varrho_{SB})] = \mathcal{L}_0(\varrho_{SB})$, meaning that measurements do not interfere with the ideal evolution.

Taylor expanding $\mathcal{U}(\tau/M) = \exp[(\tau/M)\mathcal{L}]$ in Eq. (2) and the ideal unitary superoperator $\mathcal{U}_0(\tau)$ and expanding \mathcal{L} as a sum of K -local terms yields an expression for $\varrho_{SB}(\tau) - \varrho_{SB}^0(\tau)$ as a sum of products of projectors \mathcal{P}_ϵ and Hamiltonian commutators $\mathcal{L}_{\tilde{r}_Q}(\cdot) \equiv -i[H_{\tilde{r}_Q}, \cdot]$ acting on ϱ_{SB} . By the above arguments, the projectors in each of these terms may be replaced by ζ^{jq} , where $j = 0$ if all commutators in the term are $\mathcal{L}_0 \equiv \mathcal{L}_{\tilde{r}_0}$. Invoking the triangle inequality, submultiplicativity, and the fact that $\|\varrho_{SB}\|_1 = 1$ allows the trace norm of this sum to be bounded by a linear combination of the norms of the $\mathcal{L}_{\tilde{r}_Q}$ operators, which may all then be replaced by the upper bounds $J_0 \geq \|\mathcal{L}_0\|$ and $J_1 \geq \|\mathcal{L}_{\tilde{r}_Q}\|$ or all $\tilde{r}_Q \neq \tilde{r}_0$. The resulting hypergeometric sum may be shown to equal the expression B given in Eq. (4a).

When we perform generator measurements $\bar{\mathcal{P}}_\epsilon$, each error anticommutes with at least one generator. To derive a simple but general result, we consider only the worst-case scenario of each error anticommuting with just one generator. An almost identical calculation to the one for the full stabilizer-group protocol reveals an upper bound for $D[\varrho_S(\tau), \varrho_S^0(\tau)]$ given by replacing q by 1 in Eqs. (4) and (5), since q counts the number of anticommuting stabilizer or generator elements. This completes the proof sketch. Complete proof details will be provided in Ref. [20].

We note that the generators-only bound is not as tight as the one for the full-group protocol, due to the worst-case assumption of $q = 1$ used to upper-bound terms with larger exponents which appear in the Taylor expansion discussed above. In other words, the bound on the generators-only protocol contains a sum over terms of the form $\mathcal{P}_\epsilon \mathcal{L}_{\bar{r}_Q}(\varrho) = \zeta^q \mathcal{L}_{\bar{r}_Q}(\varrho)$, where $q \in \{1, \dots, \bar{Q}\}$, all of which we have replaced for simplicity by $q = 1$. Our upper bounds are illustrated in Fig. 1. Clearly, the generators-only bound is not as close to the strong measurement limit as the full-group protocol bound. However, the former protocol requires an exponentially smaller number of measurements. If the measurement is performed, e.g., by attaching an ancilla for each measured Pauli observable (as in a typical fault-tolerant QEC implementation [12]), then this translates into an exponential savings in the number of such ancillas. Thus the two protocols exhibit a performance-resource trade-off. Next, we discuss an example.

Suppression of one-local errors.—To illustrate our general construction, we consider suppression of decoherence due to a Hamiltonian containing one-local errors on n qubits: $H_{SB} = \sum_{i=1}^n \sum_{\alpha \in \{x,y,z\}} \sigma_i^\alpha \otimes B_i^\alpha \equiv H_x + H_y + H_z$, where $J_a \equiv \|H_a\| < \infty$. This model captures the dominant errors in any implementation of quantum control or quantum computing using qubits, since any terms with higher

locality must result from three-body interactions and above. Theorem 1 guarantees first-order suppression of this H_{SB} provided we perform weak measurements over a stabilizer group of distance $d \geq 2$. We can, e.g., choose an error detection code $\mathcal{C} = [[n, n-2, 2]]$, where n is even, defined by the stabilizer generators $\bar{\mathbf{S}} = \{\bar{S}_1 = X^{\otimes n}, \bar{S}_2 = Z^{\otimes n}\}$, i.e., $\bar{Q} = 2$. The codewords are $\{|\psi_x\rangle = (|x\rangle + |\bar{x}\rangle)/\sqrt{2}\}$, where x is an even-weight binary string of length n and $x + \bar{x} = 0 \pmod{2}$. This distance $d = 2$ code is attractive since the normalizer elements are all two-local, which means that H_S is also two-local if it is constructed over these normalizer elements. Encoded single-qubit operations for \mathcal{C} are $\tilde{X}_j = \sigma_1^x \sigma_{j+1}^x$ and $\tilde{Z}_j = \sigma_{j+1}^z \sigma_n^z$, where $j = 1, \dots, n-2$. Encoded two-qubit interactions are $\tilde{X}_i \tilde{X}_j = \sigma_{i+1}^x \sigma_{j+1}^x$ and $\tilde{Z}_i \tilde{Z}_j = \sigma_{i+1}^z \sigma_{j+1}^z$. This is sufficient for universal quantum computation in both the circuit [12] and adiabatic models [21]. Thus our encoded WMQZE strategy applies in both settings (we note that our results hold even when the Hamiltonian H is time-dependent [20]). If we weakly measure the entire stabilizer group $\mathbf{S} = \{\mathbb{1}, X^{\otimes n}, Y^{\otimes n}, Z^{\otimes n}\}$, Theorem 1 implies that a state encoded into \mathcal{C} , supporting $n-2$ logical qubits, is protected by the encoded WMQZE according to Eq. (4b) with $Q = 3$ and $q = 2$. If we measure only the generators, Theorem 1 gives the same bound with ζ^2 replaced by ζ .

Conclusions.—The “traditional” QZE uses strong, projective measurements and is able only to protect an eigenstate of the operator being measured. In this work we have presented a general study of decoherence suppression via the WMQZE for arbitrary quantum states, allowing for universal quantum computation and control. By using the WMQZE to protect codewords of a stabilizer QECC, we have explicitly demonstrated that one can achieve decoherence suppression to arbitrary accuracy by increasing both the measurement strength and frequency while at the same time applying logical operators (normalizer

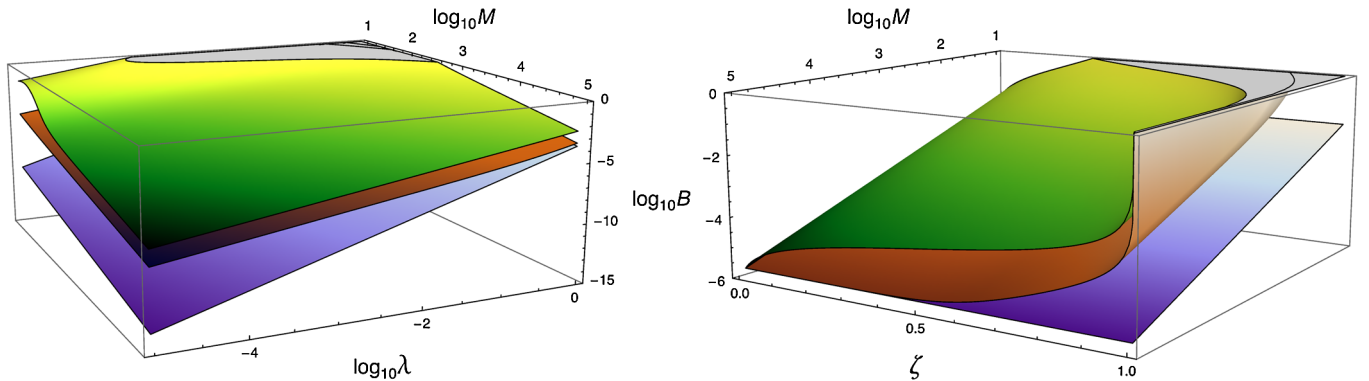


FIG. 1 (color online). Left: The upper bound B [Eq. (4a)] as a function of the number of measurements M and $\lambda \equiv J_1/J_0$, with $J_0\tau = 1$, $\bar{Q} = 4$, and $\zeta = 0.8$. Right: The same bound as a function of M and ζ , with $J_0\tau = 1$, $\bar{Q} = 4$, and $\lambda = 0.1$. In both plots the upper, middle, and lower surfaces are, respectively, the bounds for the generators-only, full stabilizer-group, and strong measurements protocols, the latter being the $\zeta \rightarrow 0$ limit of B , given in Eq. (6). The full stabilizer-group bound is tighter than the generators-only bound for all values of the parameters and is closer to the bound for the strong measurement limit.

elements) as Hamiltonians, which suffices for universality. This establishes the WMQZE as a general alternative to other open-loop quantum control methods, appropriate where measurement, rather than unitary control, is advantageous. A natural example is measurement-based quantum computation [22]. We defined two protocols, one based on measurement of the full stabilizer group and another on measurement of the generators only, and studied the trade-off between the two. The former requires exponentially more commuting measurements. However, our upper bound on its suppression of the effect of the finiteness of the measurement strength is exponentially tighter.

It would be interesting to consider whether—similarly to recent developments in dynamical decoupling theory using concatenated sequences [23] or pulse interval optimization [24–26]—WMQZE decoherence suppression can be optimized by exploiting, e.g., recursive design or nonuniform measurement intervals. Another interesting possibility is to analyze the joint effect of feedback-based quantum error correction [16] and the encoded WMQZE. Finally, an important extension would be to improve the WMQZE protocol by using techniques from fault tolerance theory [27].

D.A.L. acknowledges support from the U.S. Department of Defense and the NSF under Grants No. CHM-1037992 and No. CHM-924318. This research is partially supported by the ARO MURI Grant No. W911NF-11-1-0268.

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